

# IRPF90 : a Fortran code generator for HPC

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# Scientific codes

- Scientific codes need *speed* -> Fortran
- Fortran is a low level language -> difficult to maintain
- High-level features of Fortran 95 kill the efficiency (pointers, array syntax, etc) -> not a good solution
- Less effort in program development and good efficiency with Python/Numpy or Python/F2Py
- Other option : use Python to write low level Fortran code

# What is a scientific code?

A program is a function of its input data:

```
output = program (input)
```

A program can be represented as a **production tree** where

- The nodes are the variables
- The vertices represent the relation *needs/needed by*

Example:

```
u(x,y) = x + y + 1
v(x,y) = x + y + 2
  w(x)  = x + 3
t(x,y)  = x + y + 4
```

What is the production tree of  $t( u(d1,d2), v( u(d3,d4), w(d5) ) )$ ?

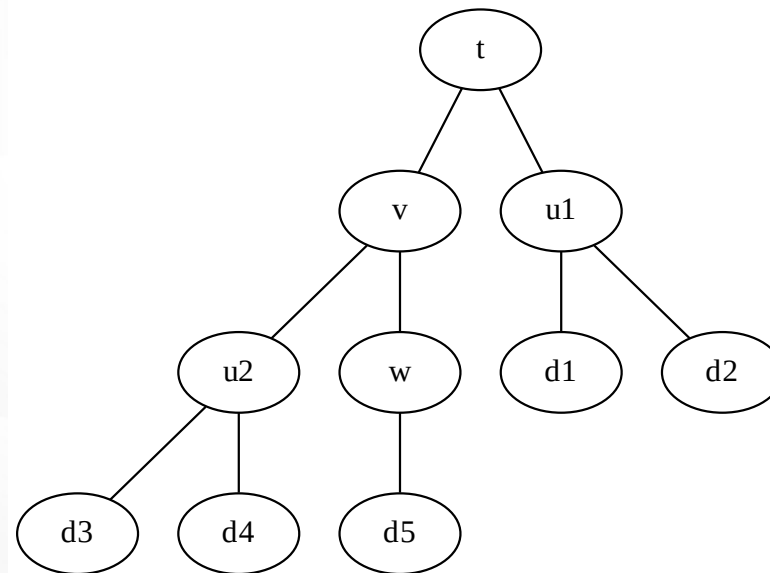
What is the production tree of  $t( u(d1,d2), v( u(d3,d4), w(d5) ) )$  )?

$$u(x, y) = x + y + 1$$

$$v(x, y) = x + y + 2$$

$$w(x) = x + 3$$

$$t(x, y) = x + y + 4$$



# Fortran way of programming

```
def fu(x,y): return x+y+1
def fv(x,y): return x+y+2
def fw(x)  : return x+3
def ft(x,y): return x+y+4
def input_data():
    # ...
    return d1, d2, d3, d4, d5

def main():
    d1,d2,d3,d4,d5 = input_data()
    u1 = fu(d1, d2)
    u2 = fu(d3, d4)
    w  = fw(d5)
    v  = fv(u2, w)
    print ft(u1, v)
```

# Difficulties

The subroutines need to be called in **the correct order**:

- The programmers needs have the knowledge of the production tree
- Production trees are usually too complex to be handled by humans
- Collaborative work is difficult : any user can alter the production tree
- Programmers may not be sure that their modification did not break some other part

# IRP way of programming

- There is only one way to build a value : by calling its *provider*
- Provider :
  - If the value is already built : return the previous value (memo function)
  - Otherwise : call the providers of the needed entities and then build the value

## Advantages:

- The provider *guarantees* that the value is valid
- Only a *local* knowledge of the production tree : the **needed** entities
- This is equivalent to calling  $t( u(d1,d2), v( u(d3,d4), w(d5) ) )$ . As there is only one set of possible parameters for each function (the needed entities are always the same), the parameters can be embedded inside the functions.

IRP : functions with Implicit Reference to Parameters

# Example

```
import sys

def irp(f):
    """All the IRP entities are represented inside a common class.
       This function generates the provider of an entity.
       f is the function that builds the entity.
    """

    # Switch the current environment to the class containing the
    # IRP entities
    locals = sys._getframe(1).f_locals

    # Name of the memo value : _f
    cache = "_" + f.__name__

    # Generic provider
```



```
def provider(self):
    # Check if self._f exists
    try:
        result = getattr(self,cache)
        print " -- Already built "+f.__name__
    except AttributeError:
        # If self._f doesn't exist, build it
        print " -> Building "+f.__name__
        result = f(self)
        setattr(self,cache,result)
        print " <- Done building "+f.__name__
    # Return self._f
    return result

# Return a class property to call the provider
return property(fget=provider)
```

```
class MyProgram(object):

    def u1(self): #  $u(x,y) = x + y + 1$ 
        return self.d1 + self.d2 + 1
    u1 = irp(u1)

    def u2(self): #  $u(x,y) = x + y + 1$ 
        return self.d3 + self.d4 + 1
    u2 = irp(u2)

    def v(self): #  $v(x,y) = x + y + 2$ 
        return self.u2 + self.w + 2
    v = irp(v)

    def w(self): #  $w(x) = x + 3$ 
        return self.d5 + 3
    w = irp(w)
```

```
def t(self): # t(x,y) = x + y + 4
    return self.u1 + self.v + 4
t = irp(t)
```

```
def d(self):
    return range(1,6)
d = irp(d)
d1 = property(lambda self: self.d[0])
d2 = property(lambda self: self.d[1])
d3 = property(lambda self: self.d[2])
d4 = property(lambda self: self.d[3])
d5 = property(lambda self: self.d[4])
```

```
def main():
    p = MyProgram()
    print "u1 : " ; print p.u1
    print "t  : " ; print p.t
```

```
main()
```

```
u1 :          #          t
-> Building u1 #          /          \
-> Building d  #          u1          v
<- Done building d # / |          | \
-- Already built d # d1 d2          u2 w
<- Done building u1 #
4                   #          /          \          \
t :                  #          d3          d4          d5
-> Building t
-- Already built u1
-> Building v
-> Building u2
-- Already built d
-- Already built d
<- Done building u2
-> Building w
```

```
-- Already built d  
<- Done building w  
<- Done building v  
<- Done building t
```

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# How to do this with Fortran?

Fortran doesn't have exceptions, introspection, properties, etc.

Solution IRPF90:

- Add a few keywords to Fortran
- Use Python to read the code
- Python builds the dependence tree
- Python writes the missing Fortran source lines to handle the tree

## Example with IRPF90

```
BEGIN_PROVIDER [ integer, t ]  
  t = u1+v+4  
END_PROVIDER  
  
BEGIN_PROVIDER [ integer, w ]  
  w = d5+3
```

```
END_PROVIDER
```

```
BEGIN_PROVIDER [ integer, v ]
```

```
  v = u2+w+2
```

```
END_PROVIDER
```

```
BEGIN_PROVIDER [ integer, u1 ]
```

```
  integer :: fu
```

```
  ! u1 = fu(d1,d2)
```

```
  u1 = d1+d2+1
```

```
END_PROVIDER
```

```
BEGIN_PROVIDER [ integer, u2 ]
```

```
  integer :: fu
```

```
  ! u2 = fu(d3,d4)
```

```
  u2 = d3+d4+1
```

```
  ASSERT (u2 > d3)
```

```
END_PROVIDER
```

```
integer function fu(x,y)
  integer :: x,y
  fu = x+y+1
end function
```

```
program irp_example
  print *, 't = ', t
end
```



# Features

## Arrays

```
BEGIN_PROVIDER [ double precision, A, (dim1, 3) ]  
  . . .  
END_PROVIDER
```

- Allocation of IRP arrays done automatically
- Dimensioning variables can be IRP entities, provided before the memory allocation

# Documentation

Every subroutine/function/provider can have a documentation section:

```
BEGIN_PROVIDER [ double precision, Fock_matrix_beta_mo, (mo_tot_num_align,mo_tot_num) ]
  implicit none
  BEGIN_DOC
  ! Fock matrix on the MO basis
  END_DOC
  ...
END_PROVIDER
```

```
$ irpman fock_matrix_beta_mo
```

```

IRPF90 entities(1)                fock_matrix_beta_mo                IRPF90 entities(1)

Declaration
    double precision, allocatable :: fock_matrix_beta_mo    (mo_tot_num_align,mo_tot_num)

Description
    Fock matrix on the MO basis

File
    Fock_matrix.irp.f

Needs
    ao_num
    fock_matrix_alpha_ao
    mo_coef
    mo_tot_num
    mo_tot_num_align

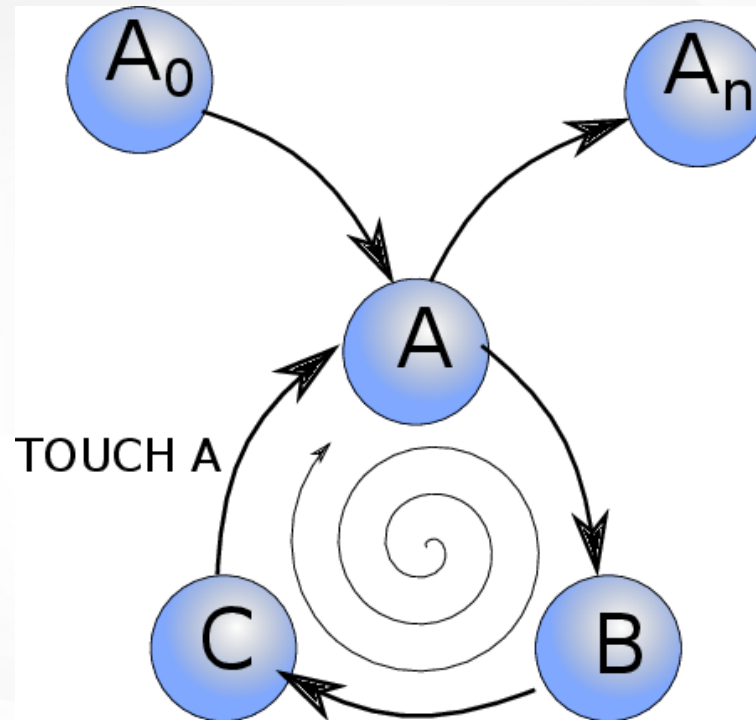
Needed by
    fock_matrix_mo

IRPF90 entities                fock_matrix_beta_mo                IRPF90 entities(1)

```

# Iterative processes

Iterative processes may involve cyclic dependencies:



TOUCH A : A is valid, but everything that needs A is invalidated

# Meta-programming

```
BEGIN_SHELL [ /bin/bash ]
  echo print *, \'Compiled by `whoami` on `date`\`\'
  echo print *, \'$FC $FCFLAGS\'
  echo print *, \'$IRPF90\'
END_SHELL
```

```
BEGIN_SHELL [ /usr/bin/python ]
for i in range(10):
  print """
    double precision function times_%d(x)
      double precision, intent(in) :: x
      times_%d = x*%d
    end
  """%locals()
END_SHELL
```

# Many Other features

- Color highlighting in Vi
- Generation of tags to navigate in the code
- Variables can be declared *anywhere*
- Dependencies are known by IRPF90 -> Makefiles are built automatically
- No problem using external libraries (MKL, MPI, etc)
- Compatible with OpenMP
- Support for Coarray Fortran (distributed parallelism)
- Codelet generation for code optimization
- Generation of Intel Fortran compiler directives to align arrays
- Generated code is **very** efficient (960 Tflops/s on Curie in 2011 with QMC=Chem)
- etc...